

# Garbled Circuits: Intro to MPC

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# Multiparty Computation

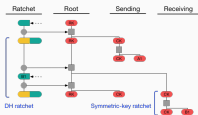
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# Cryptography

## Secure Communication

- AES, RSA, etc.
- WhatsApp, HTTPS



**Figure 1:** Double Ratchet:  
Signal

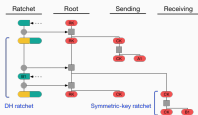


**Figure 2:** TLS: Cloudflare

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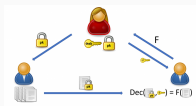
**Figure 1:** Double Ratchet: Signal



**Figure 2:** TLS: Cloudflare

## Secure **Computation**

- Modern day constructions
- MPC, FE, FHE, ZK, and more
- Secure Voting and Auctions, Cryptocurrency, Secure ML



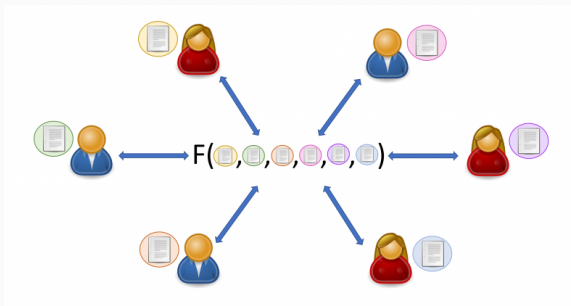
**Figure 3:** MPC: COSIC

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**Figure 4:** Secure MPC: Cosic

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- Hiding auction bids on a smart contract

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- Lets build it from the ground up!

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- First two party protocols introduced by Yao - Garbled Circuits
- Lets build it from the ground up!
- Compute function between two parties  $P_1$  and  $P_2$
- Both  $P_1$  and  $P_2$  are honest

# Garbled Circuits

## Some Notation

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Let  $Y = \{y_0, y_1, \dots\}$  be the set of possible  $y$ .
- $P_1$  is the party choosing  $x$ ,  $P_2$  is the party choosing  $y$
- $P_1, P_2$  want to compute  $\mathcal{F}(x_a, y_b)$  without revealing  $a, b$ .

- Given a function  $\mathcal{F}(x, y)$ , how do we evaluate  $\mathcal{F}(x_a, y_b)$ ?

# Garbled Circuits

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- The function  $\mathcal{F}$  is just a table!

**Table 1:** A function with  $X = \{x_0, x_1\}$ ,  $Y = \{y_0, y_1\}$

	$y_0$	$y_1$
$x_0$	$\mathcal{F}(x_0, y_0)$	$\mathcal{F}(x_0, y_1)$
$x_1$	$\mathcal{F}(x_1, y_0)$	$\mathcal{F}(x_1, y_1)$

## Garbled Circuits

- $P_1$  assigns key  $k_i^x$  to  $x_i$ , and assigns key  $k_i^y$  to  $y_i$ .

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**Table 2:** A Garbled Function  $\tilde{\mathcal{F}}$

	$y_0$	$y_1$
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- We can send the encrypted values to  $P_2$
- If  $P_2$  has the two keys associated to  $x_a, y_b$ , they can get the output!

## Garbled Circuits

- Problem:  $P_2$  knows which row/column is based on which value



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- Pointer value tells  $P_2$  which row/column is correct
- $p_0^x, p_1^x = 1 - p_0^x$

## Garbled Circuits

- $P_1$  sends  $P_2$  the garbled  $\tilde{\mathcal{F}}$
- $P_1$  sends  $P_2$  the key  $k_a^x$  and  $p_a^x$
- How does  $P_2$  get  $k_b^y$  without  $P_1$  learning  $b$ ?

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- Alice has  $k_0^y$  and  $k_1^y$
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# Oblivious Transfer

- Introducing a new protocol: Oblivious Transfer!
- Alice has  $k_0^y$  and  $k_1^y$
- Bob has a bit  $b$
- Alice learns nothing, and Bob learns  $k_b^y$

# Oblivious Transfer

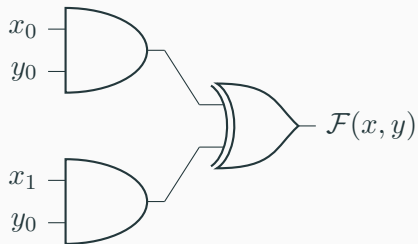
- Oblivious transfer exists
- Probably ask me later about how to construct OT
- Relies on some more cryptography concepts not taught

- Back to the garbled table
- $P_1$  sends  $\tilde{\mathcal{F}}, k_a^x, p_a^x$
- $P_1$  and  $P_2$  do Oblivious Transfer to send  $k_b^y, p_b^y$  to  $P_2$

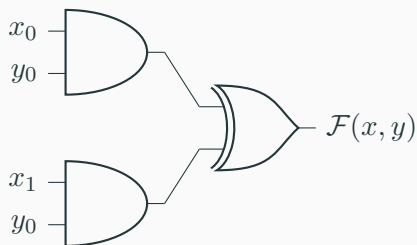


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- With both pointers,  $P_2$  can find the location in the table
- With both keys,  $P_2$  can decrypt that row, to get the output



**Figure 5:** Example Circuit



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- Lookup table?? Exponential

**Table 4:** A Garbled Gate  $\tilde{\mathcal{F}}$

	$y_0$	$y_1$
$x_1$	$\text{Enc}(k_1^x \oplus k_0^y, k_{\mathcal{F}(x_1, y_0)}^{\text{out}})$	$\text{Enc}(k_1^x \oplus k_1^y, k_{\mathcal{F}(x_1, y_1)}^{\text{out}})$
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## Garbled Circuits

- Assign pair of keys (and pointers) **per wire**
- Can evaluate circuits gate by gate!

## More Fun Things

- This is just the surface: many more questions
- How can we do this with multiple parties?
- Can we improve communication efficiency?



## More Fun Things

- Very useful in the world of secure computation
- For example, say the circuit is a circuit which takes in an ML classifier and an image, and outputs the classification
- Allows doctors to provide medical images / data to ML models without violating HIPAA

Bonus: Funny card trick (related to MPC)!!

$(C, C)C(C, C)$

Both say yes:

C C C C C

One or both says no:

C C C C C